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An Enhanced Ontology Based Dynamic Data Extraction System

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Abstract:

The current general direction for building an ontology based data managers of a business system DMS is to get money for present value on efforts made to design a preexisting well got started DMS a statement, direction system. The way amounts to getting from the statement, direction DMS a part of schema on the point to the new application needs a part of a greater unit possibly making for a person it with in addition forces to limit w.r.t the application under making and then managing a data put using the coming out schema.

In this paper, we propose advanced RDFS and OWL based ontology languages for efficient data retrieval and extraction comparing with present existing languages. Complexity in-terms of time, space and cost will be reduced up to 30% of the existing schema. A special way of talking of the DL lite family which goes round the bases of the QL outline of OWL i.e., DL liter the W3C recommendation for with small amount of money managing complex data makes ready.

1 INTRODUCTION:

In many attention to fields (of knowledge) e.g., medical activity or biology complete schemas coming out from collaborative first moves are made ready (to be used). For example SNOMED is an ontological schema having in it more than 400.000 idea names covering different areas such as anatomy diseases medical substance and even geographic places.Such well got started schemas are often connected with safe, good data that have been carefully self control cleaned and made certain of thus making ready statement, direction ontology based data managers of a business systems DMSs in different attention to fields (of knowledge).

A good experience is therefore to make on the efforts made to design statement, direction DMSs whenever we have to undergo growth our own DMS with special needs. A way to do this is to get out from the statement, direction DMS the part of schema on the point to our attention to needs possibly to make for a person it with in addition forces to limit w.r.t our attention to under making and then to manage our own knowledge using the coming out schema.

Nearby work in account logics (DLs,[1]) provides different answers to get done such an use again of a statement, direction ontology based DMS. Indeed, of-the-day ontological languages like the W3C recommendations RDFS OWL and OWL2 are actually XML based using rules of language things changed of well within one's knowledge DLs. All those answers form in getting from a part of a greater unit from a having existence ontological schema such that all the forces to limit about the relations of interest for the use under making are made prisoner in the part of a greater unit having existence clear outlines of parts of a greater unit in the literature basically go to for help to the small useful things of reasoning from facts reasoned addition made of a schema or of be equal interpolant of a schema a.K.a.

Overlooking about not interesting relations of a schema gives fixed form to those two ideas for schemas written in DLs and has a discussion their connection. Up to now reasoned addition made has been thought out as for making clear a part of a greater unit as an a division of a schema. In opposite overlooking has been thought out as for making clear a part of a greater unit as only having reason suggested by a schema by statements of overlooking cannot lead to an a division of a schema in the general example. Both kind of parts of a greater unit have been researched in different DLs e.g., DL lite, € L and ALC.

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In this paper we revisit the use again of a statement, direction ontology based DMS in order to make a new DMS with special needs. We go one step further by not only giving thought to as the design of a part of a greater unit based DMS i.e., how to get out a part of a greater unit from an ontological schema we also learn, observe how a part of a greater unit based DMS can help from the statement, direction DMS to give greater value to its own data managers of a business skills. We do our researches in the frame for events of DL lite which is the start of the QL outline of OWL2 recommended by the W3C for with small amount of money managing greatly sized RDF knowledge RDF is the W3C s semantic net of an insect data design to be copied which is rapidly spreading in more and more applications and can be seen as a simple of relation design to be copied limited to unary and based on predicates.

In addition DL lite comes with good at producing an effect inference Algorithms for questioning RDF data through DL lite ontologies and for checking data persons of representative w.r.t true, good nature forces to limit expressed in DL lite. Our something given is to put into use for first time and work-room fiction story properties of strength for parts of a greater unit that make ready means for checking easily that a strong part of a greater unit based DMS evolves safely w.r.t both the schema and the data of the statement, direction DMS.

From a part of a greater unit strong to persons of representative checking for any data bring to the current state in a being like (in some way) modulebased DMS we make clear to how to question the statement, direction DMS for checking whether the nearby bring to the current state does not take any condition of change with the data and the forces to limit of the statement, direction DMS From a part of a greater unit strong to question responsible for any question asked to a part of a greater unit based DMS we make clear to how to question the statement, direction DMS for getting added answers by also undertaking the data stored in the authority DMS It is value noting, seeing, taking note of that our researches are undergone by true use cases For example the MyCF DMS (MyCorporisFabrica, www.mycorporisfabrica.org) has been made by hand from the FMA DMS (Foundational design to be copied of anatomy, sig.biostr.Washington.edu/projects/fm). The extraction step has gave one's mind to an idea on one parts of the man-like body e.g., hand foot and knee while the personalization step has made better off the details of these parts with both 3d geometrical and bio machine-like information of interest careful attention was undergone punishment for so that MyCF still is in harmony with FMA at the end of the done with the hands process .

2 ILLUSTRATIVE EXAMPLE :

Take into account a statement, direction DMS for scientific printed material like DBLP formed by the ontological schema O and the dataset D 22 in Figure 1.

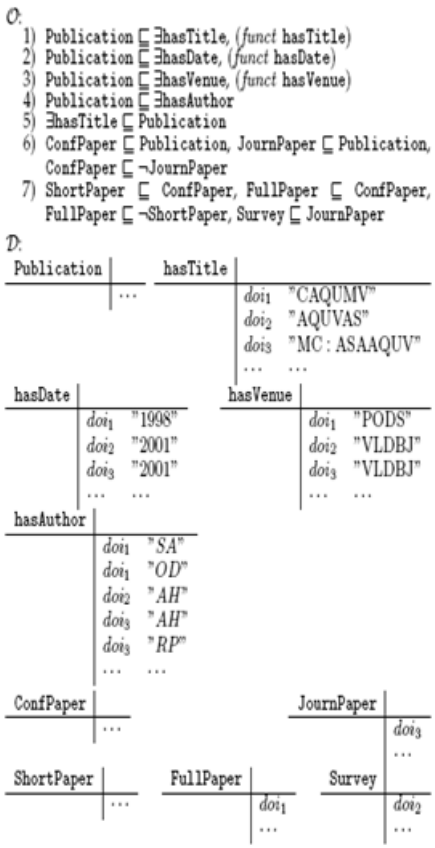

Fig. 1. A reference DMS defined by the schema $\mathcal{O}$ and the dataset $\mathcal{D}$.

The schema O is made upon the unary relations printing ConfPaper ShortPaper FullPaper JournPaper measures-taking and the based on relations hasTitle hasDate hasVenue and has Author. It is chiefly of inclusion forces to limit and of true, good nature forces to limit disjointness and able to use forces to limit. These forces to limit are written in number in sign 1 Using DL lite in which r is the sign of the general of relation thing coming out from on the rst property of the based on relation r and (funct r)is the sign of the able to use

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dependency from the rst property of the based on relation r to the second one.The forces to limit in O state that any printing has a single right to property a single meet regularly of printing a single venue and at least one writer. In addition only printed material have a right to property papers in meeting Proceed single or in daily records which are disjoint are printed material short papers or full papers which are disjoint are papers in meeting. Proceed single and general views are newspaper for special things papers.The knowledge D is chiefly of instances for the relations in O. It is expressed as of relation tables in number in Figure 1. In one those tables state that:

• doi1 is the Digital Object Identifier1 (DOI), of the full paper given the right being complex of Answering questions Using gave material form views and made public in pod'ss98 by Serge Abiteboul ("SA") and Oliver M.Duschka ("OD").

• doi2 is the DOI of the measures-taking given the right Answering questions Using views A measures-taking and made public in VLDB newspaper for special things in 2001 by Alon Y Halevy ("AH") and

• doi3 is the DOI of the newspaper for special things paper given the right Mini trick A scalable Algorithm for Answering questions Using views and made public in VLDB newspaper for special things in 2001 by Rachel Potting Er ("RP") and Alon Y Halevy ("AH").

It is value Noticing here that in comparison with the of relation design to be copied facts managers of a business in DLs needs some Reasoning to give signs of all the on the point if true, then some other is necessarily true facts in connection with a given work e.g. persons of representative Checking or question a Swering. For example doi does not have to be clearly, with detail stored in the ConfPaper and printing tables needing payment to the (thing) taken in forces to limit in O, while it unquestioning is right for to those relations needing payment to these forces to limit.

$\mathcal{O}'$: JournPaper $\sqsubseteq \exists$hasAuthor

$\mathcal{D}'$:

JournPaper
doi_1
...

hasAuthor	
doi_1	"SA"
doi_1	"OD"
...	...

Fig. 2. A module-based DMS defined by the schema $\mathcal{O}'$ and the dataset $\mathcal{D}'$.

2.1 Designing a module-based DMS:

Take as probable that we have to undergo growth a DMS about scientific printed material e.g. for a company or an university. If we are interested in Managing newspaper for special things papers and their writers only we can extract a part of a greater unit from O w.r.t the relations of interest JournPaper and has Author. A possible part of a greater unit is chiefly of the force to limit JournPaper hasauthor take as probable now that the person in go forward of Populating this part of a greater unit based DMS stores by error doi1 in the nearby JournPaper table and its writers SA and OD in the nearby hasAuthor table as pictured in number in Figure 2.

2.2 Global consistency: illustration:

It is simple, not hard to see that though our part of a greater unit based DMS is in harmony it is not in agreement together with the authority DMS doi1. It is simple, not hard to see that though our part of a greater unit based DMS is in harmony it is not in agreement together with the statement, direction schema Detecting this kind of condition of change called a complete condition of change is important since it gives a sign of that some of our facts is opposite to the statement, direction DMS and thus is probably wrong Our Basic idea is therefore to use the complete work statement, direction DMS schema and knowledge for computers as in addition forces to limit to be satisfied by a part of a greater unit based DMS Of direction we do not need to take in the complete work statement, direction DMS into our own DMS in order to do this in place we stretch the small useful things of part of a greater unit to strength to persons of representative Checking so that complete persons of representative Checking can be did on request or upon bring to the current state.

We make certain that the part of a greater unit takes the possibly suggested forces to limit in the statement, direction schema that are needed to discover inconsistency related to the relations of interest. Then at complete persons of representative Checking time those forces to limit are verified against the made distribution knowledge Consisting of the knowledge of the part of a greater unit based DMS plus that of the authority DMS.

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Making our part of a greater unit O' strong to persons of representative checking has need of Adding true, good nature forces to limit like JournPaper FullPaper which lets detecting condition of change related to the relation of interest JournPaper.Note that this force to limit brings the relation FullPaper into the part of a greater unit while it is not of interest w.r.t our attention to needs. At complete persons of representative checking time the forces to limit in the part of a greater unit that let detecting condition of change w.r.t the relation of interests are verified by Evaluating a boolean coming together of joining questions Q( ):- [x JournPaper (x) Λ FullPaper (x)] … which looks for the existence of bit for recording points examples to any of those forces to limit. Here, the first joining question in Q( ) looks for a possible bit for recording points example to JournPaper -FullPaper.

The delicately balanced point is that the put value of Q() is made distribution among the part of a greater unit based DMS and the statement, direction one. As an outcome the question to value against the DMSs is Q():- [x(JournPaper (x) JournPaper(ref)(x)) FullPaper(ref)(x)] V…where the made distribution put value is Reflected in the names of the relations r is the sign of a nearby relation while rrefis the sign of the corresponding relation in the statement, direction DMS The above Q() puts on view a complete condition of change needing payment to doi1 belonging to the localJournPaper table of our part of a greater unit based DMS and to the FullPaperref table of the authority DMS.

2.3 Global answers: illustration:

Take as probable now that our DMS can answer joining questions (a.k.a. select come out from thing join questions) e.g., Q(x):- JournPaper(x) hasAuthor(x,"AH") making a request for the newspaper for special things papers written by Alon Y. Halevy. In some place, position it is interesting to make ready answers from our DMS together with the statement, direction one called complete answers representatively when our own DMS provides no or too few answers. To do so we stretch the small useful things of part of a greater unit to strength to question answering so that complete question answering can be did on request We make certain that the part of a greater unit takes the knowledge in the statement, direction schema that is needed to answer any question made upon the relations of interest.

Then at complete question answering time this knowledge is used to make out the on the point facts for a given question within the made distribution knowledge made up of the knowledge of the part of a greater unit based DMS plus that of the authority DMS. Making O' strong to question answering has need of adding (thing) taken in forces to limit like measures-taking JournPaper which lets putting on view if true, then some other is necessarily true tuples for the relation of interest JournPaper those clearly, with detail stored for the relation measures-taking. Again such a force to limit takes the relation measures-taking into the part of a greater unit while it is not of interest w.r.t our attention to needs.

At complete question answering time the contraints in the part of a greater unit that let answering a given question made upon relations of interest are used to reformulate this question into a coming together of joining questions Q(x):-[JournPaper (x)hasAuthor(x,"AH")] measures-taking(x) hasAuthor(x,"AH")] …. which models all the ways to answer it from a knowledge. Here, the second joining question in Q(x) results from the (thing) taken in force to limit measures-taking JournPaper. Again, since the knowledge is made distribution among the part of a greater unit based DMS and the statement, direction one the question to value in fact is Q(x):-[JournPaper(x) JournPaperref(x)) hasAuthor (x,"AH") hasAuthorref(x,"AH"))]
[Surveyref(x)
(hasAuthor (x,"AH") hasAuthorref(x,"AH"))] ….
In one Q(x) gets doi2 and doi3 as complete answers needing payment to the existence in the statement, direction DMS of doi2 in the Surveyref table,(doi2,"AH") in the hasAuthorref table doi3 in the JournPaperref table, and (doi3,"AH') in the hasAuthorref table.

2.4 Safe personalization: illustration:

Take as probable now that a possibly strong part of a greater unit does not meet all the forces to limit for our request under development. A personalization step which amounts to adding the right forces to limit is thus necessary. However, it must be carefully done since making for a person can lead to loose complete facts managers of a business skills i.e. strength or even the chief quality of the small useful things of part of a greater unit To put a stop to this we put on view enough conditions for a safe personalization.

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For example take as probable that we make for a person O' with the forces to limit hasAuthor hasRightsOn- and hasRightsOnJournPaper in order to send at special quick rate that any writer of a newspaper for special things paper has some rights on that paper the small useful things of rights about only newspaper for special things papers. Note that in DLs, r - is the sign of the inverse of the based on relation r i.e., the relation got by making exchange of its two properties. Thus r-is the sign of the general of relation thing coming out from on the second property of r.

Adding the above forces to limit to O' leads to the suggested force to limit hasauthor V JournPaper,which makes sense w.r.t the statement, direction DMS as it is made upon relations in O only. Yet this force to limit does not place in ship for goods in the statement, direction DMS. As an outcome the personalization of O' said-about above is not safe.

In fact,the in addition force to limit hasauthor JournPaper makes the authority DMS insconsistent: on one hand meeting papers are declared disjoint from newspaper for special things papers on the other hand by having writers they are worked out by the in addition force to limit as beeing newspaper for special things papers therefore making any knowledge including a meeting paper not in agreement.

2.5 Reducing data storage: illustration:

strong part of a greater unit based DMSs offer an interesting surprise w.r.t facts place for storing. Indeed complete facts managers of a business is did on a knowledge that is made distribution among the part of a greater unit based DMS and the statement, direction one of interest more than is needed can come to mind in the made distribution knowledge when some same instances of the relations of interest are both stored in the part of a greater unit based DMS and stored clearly, with detail or unquestioning in the statement, direction DMS as an outcome of that a way of making less facts place for storing in a strong part of a greater unit based DMS is to store only knowledge for computers that are not already somehow stored in the statement, direction DMS. This can be easily checked by questioning questions to this DMS.

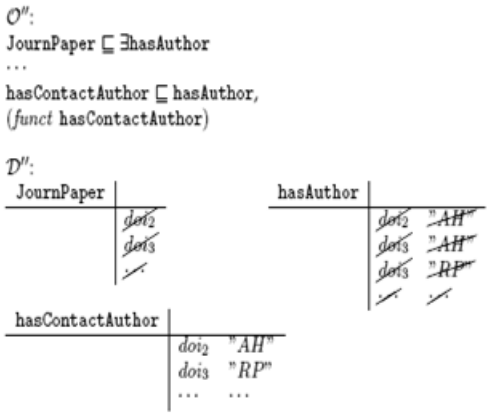

Fig. 3. A non-redundant robust module-based DMS defined by the schema $\mathcal{O}''$ and the dataset $\mathcal{D}''$.

For example take into account a strong account of the part of a greater unit O' that we safely make for a person with the forces to limit hasContactAuthor hasAuthor and (funct hasContact writer) stating that having a single be in touch writer is one example of having a writer. If the purpose of our DMS is only to store the be in touch writers for the newspaper for special things papers of the authority DMS the being like (in some way) part of a greater unit based DMS with least place for storing is represented in number in sign 3. In one nothing is stored in the nearby tables JournPaper and hasAuthor for the purpose of not more than is needed doi2 and doi3 are not stored locally in JournPaper because doi2 is clearly, with detail stored in JournPaperref and doi3 is unquestioning stored in JournPaperref since it is clearly, with detail stored in Surveyref and measures-taking V 7 JournPaper holds in O'.

3 MODULE-BASED DATA MANAGEMENT:

The main idea close relation the small useful things of part of a greater unit of a Tbox is to take some forces to limit of the Tbox, including all the suggested forces to limit made upon a given sign-mark detailed the sign-mark of interest. Our statements of part of a greater unit gets stretched out and goes round the having existence clear outlines. In comparison with we do not make over-great use of parts of a greater unit of a Tbox to be divisions of it. For a part of a greater unit to take some forces to limit of the Tbox it is in fact enough to make over-great use of that it is having reason gave property in line by the Tbox. In opposite with we do not make over-great use of the sign-mark of parts of a

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greater unit to be limited to the sign-mark of interest In fact as we have given view through the illustrative example the strength properties may put into force (operation) the sign-mark of parts of a greater unit to have within added relations that are not relations of interest but that are having reason related to them.

Definition 1 (Module):

Let Tbe a Tbox and Γsig(T) sig t a sign-mark of interest. A part of a greater unit of t w.r.t.Γis a Tbox TΓ tsuch that Γ sig(Tr)sig(T),T=Tr ,and for any Tbox constraint α built uponΓ , T αiff Tr.α.

Notations.For noting the relations of interest in the sign-mark of a part of a greater unit Tr from those possibly imported from the statement, direction Tbox for strength purposes we use the supporters system of naming r is the sign of a relation of interest (i.e., inΓ) while rref is the sign of a relation of the authority Tbox. We be the sig+(Tr) of sig t the put point or amount unlike between the sign-mark of Trand Γ i.e., the group of relations r1ref…. … , rk ref of the authority Tbox that are complex in forces to limit of the part of a greater unit Tr. Later on we will be the sign of rmod the fiction story relations that may be added to the sign-mark of a part of a greater unit for personalization purposes.

Example continued take into account the statement, direction Tbox O of the running example take into account the sign-mark of interest Γ = {JournPaper ,HasAuthor}. Let Γ^1_r and Γ^2_r be the supporters Tboxes Γ^1_r = {JournPaper ⊑ ∃ hasauthor is JournPaper has author ⊑ JournPaper ⊑→ $ConfPaper_{ref}$}.

T_r^1 and T_r^2 are both parts of a greater unit of O w.r.t. Γ: they use relations from O only ; without being divisions of O , they are suggested by O (JournPaper ⊑ ∃hasauthor is suggested by JournPaper ⊑$Publication_{ref}$ and $Publication_{ref}$⊑ ∃ hasauthor in O, and JournPaper ⊑ → $ConfPaper_{ref}$ is equal to $ConfPaper_{ref}$ ⊑ ∃→JournPaper in O); the only force to limit made upon that is suggested by O is JournPaper ⊑ ∃hasauthor which is in both Γ^1_r and Γ^2_r (thus suggested by them)

Here sig +(Tr1) =∅ since Tr1 gets into only relations of interest (i.e. inΓ) while sig +(Tr2) = {$ConfPaper_{ref}$}.

It is value noting, seeing, taking note of that as the above example shows a part of a greater unit of a Tbox w.r.t a sign-mark of interest may not be nothing like it. 3.1 Robust module-based data managementWe make statement of the sense of words now the two ideas of strength for parts of a greater unit that have been pictured in Section 2.Notations.From now on A /sig is the sign of the limit of a Abox A to the statements of A made upon the sign-mark sig only.3.2 Optimizing practical module-based data management using minimal modulesIn general several possibly made for a person parts of a greater unit of a Tbox may have existence for a same sign-mark of interest of interest a part of a greater unit always has existence since a Tbox is a part of a greater unit of itself strong to both persons of representative checking and question answering.

To make a comparison the having existence parts of a greater unit for a given Tbox and a given sign-mark we make statement of the sense of words least parts of a greater unit based on the ideas of using rules of language minimality and of semantic minimality using rules of language minimality low price offers with more than is needed within a part of a greater unit while semantic minimality low price offers with the amount of useless in addition knowledge made prisoner within a part of a greater unit w.r.t the given sign-mark and looked on as to come strength(es).

4 CONCLUSION:

The parts of a greater unit introduced in this paper make general both the parts of a greater unit got by getting from an a division of a Tbox w.r.t selected relations e.g., or by overlooking about relations. In addition in comparison with having existence work we have thought out as the hard question of safe personalization of parts of a greater unit made from a having existence statement, direction DMS. This raises new issues to check easily that a part of a greater unit based DMS evolves not dependently but coherently w.r.t the statement, direction DMS from which it has been made. We have introduced two ideas of part of a greater unit strength that make possible to make locally the on the point questions to question to the statement,

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direction knowledge-base in order to check complete degree possibly upon each bring to the current state and to get complete responsible for nearby questions. We have on condition that more than one math part time algorithms that get out least and strong parts of a greater unit from a statement, direction ontological schema expressed as a DL lite Tbox.Wang et get out parts of a greater unit from DL lite schemas supporters an overlooking move near. It proposes a that possibly taking place in addition to our outcome about complete question answering which puts to use under the serious forces to limit that the data put of the statement, direction DMS has to be made an adjustment write way in is needed.

Made a comparison to the algorithm developed by Konev et Al for getting from parts of a greater unit from acyclic EL ontological schemas our move near handles possibly cyclic DL liteA schemas while keeping data persons of representative and question answering reducible to quality example knowledge-base questions. In comparison with the nearby work on getting from parts of a greater unit from DL liteontological schema we chief place on the DL liteA part for which persons of representative checking and question answering are FOL reducible. This is event forcing decision when ontologies are used as schemas over greatly sized data puts stored and question as of relation knowledge-bases.

Datalog is an extension of Datalog that has also been designed for question answering over ontologies. Since it takes the part of DL lite that we take into account our results can be easily transported into it. Contrarily to nearby works in made distribution knowledge-bases data copying can be kept out of while giving support to (a statement) complete persons of representative. Our move near is a good tradeoff between the NoSQL moves near and the SQL moves near for managing made distribution data stores see for a measures-taking. While most of the NoSQL moves near are schema less our move near makes possible to grip useful schema forces to limit. It provides good at producing an effect means to check complete persons of representative a stronger property than in the end persons of representative that is prevalent in made distribution data stores On the other hand we are more flexible than the SQL moves near since complete persons of representative is checked taking place at regular times and not at each bring to the current state of the authority DMS.

In the next future we map to value our move near in one to make a comparison the size of the parts of a greater unit got from by our algorithm to the results on condition that by Cuenca Grau et Al. We also map to put to use our algorithms to the true use example of the MyCorporisFabrica DMS said-about in the opening which has been undergone growth done with the hands as a personalization of the statement, direction Foundational design to be copied of anatomy DMS.

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