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Improving Performance of Secure Data Transmission in Communication Networks Using Physical Implementation of AES

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ABSTRACT:

The importance of cryptography applied to security in electronic data transactions has acquired an essential relevance during the last few years. A proposed FPGAbased implementation of the Advanced Encryption Standard (AES) algorithm is presented in this paper. The design has been coded by Very high speed integrated circuit Hardware Descriptive Language. All the results are synthesized and simulated using Xilinx ISE and ModelSim software respectively.

This implementation is compared with other works to show the efficiency. The design uses an iterative looping approach with block and key size of 128 bits, lookup table implementation of S-box. This gives low complexity architecture and easily achieves low latency as well as high throughput. Simulation results, performance results are presented and compared with previous reported designs.

1. INTRODUCTION:

AES is short for Advanced Encryption Standard and is a United States encryption standard defined in Federal Information Processing Standard (FIPS) 192, published in November 2001. It was ratified as a federal standard in May 2002. AES is the most recent of the four current algorithms approved for federal us in the United States.

One should not compare AES with RSA, another standard algorithm, as RSA is a different category of algorithm. Bulk encryption of information itself is seldom performed with RSA.RSA is used to transfer other encryption keys for use by AES for example, and for digital signatures. AES is a symmetric encryption algorithm processing data in block of 128 bits.

A bit can take the values zero and one, in effect a binary digit with two possible values as opposed to decimal digits, which can take one of 10 values. Under the influence of a key, a 128-bit block is encrypted by transforming it in a unique way into a new block of the same size. AES is symmetric since the same key is used for encryption and the reverse transformation, decryption. The only secret necessary to keep for security is the key.

AES may configured to use different key-lengths, the standard defines 3 lengths and the resulting algorithms are named AES-128, AES-192 and AES-256 respectively to indicate the length in bits of the key. Each additional bit in the key effectively doubles the strength of the algorithm, when defined as the time necessary for an attacker to stage a brute force attack, i.e. an exhaustive search of all possible key combinations in order to find the right one.

2. METHODOLOGY: 2.1 Some background on AES:

In 1997 the US National Institute of Standards and Technology put out a call for candidates for a replacement for the ageing Data Encryption Standard, DES. 15 candidates were accepted for further consideration, and after a fully public process and three open international conferences, the number of candidates was reduced to five.

In February 2001, the final candidate was announced and comments were solicited. 21 organizations and individuals submitted comments. None had any reservations about the suggested algorithm. AES is founded on solid and well-published mathematical ground, and appears to resist all known attacks well. There's a strong indication that in fact no back-door or known weakness exists since it has



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been published for a long time, has been the subject of intense scrutiny by researchers all over the world, and such enormous amounts of economic value and information is already successfully protected by AES. There are no unknown factors in its design, and it was developed by Belgian researchers in Belgium therefore voiding the conspiracy theories sometimes voiced concerning an encryption standard developed by a United States government agency. A strong encryption algorithm need only meet only single main criteria: There must be no way to find the unencrypted clear text if the key is unknown, except brute force, i.e. to try all possible keys until the right one is found.

A secondary criterion must also be met:

The number of possible keys must be so large that it is computationally infeasible to actually stage a successful brute force attack in short enough a time. The older standard, DES or Data Encryption Standard, meets the first criterion, but no longer the secondary one – computer speeds have caught up with it, or soon will. AES meets both criteria in all of its variants: AES-128, AES-192 and AES-256.

2.2Encryption must be done properly:

AES may, as all algorithms, be used in different ways to perform encryption. Different methods are suitable for different situations. It is vital that the correct method is applied in the correct manner for each and every situation, or the result may well be insecure even if AES as such is secure. It is very easy to implement a system using AES as its encryption algorithm, but much more skill and experience is required to do it in the right way for a given situation. No more than a hammer and a saw will make anyone a good carpenter, will AES make a system secure by itself. To describe exactly how to apply AES for varying purposes is very much out of scope for this short introduction.

2.3 Strong keys:

Encryption with AES is based on a secret key with 128, 192 or 256 bits. But if the key is easy to guess it doesn't matter if AES is secure, so it is as critically vital to use good and strong keys as it is to apply AES properly. Creating1 good and strong keys is a surprisingly difficult problem and requires careful design when done with a computer.

The challenge is that computers are notoriously deterministic, but what is required of a good and strong key is the opposite – unpredictability and randomness. Keys derived into a fixed length suitable for the encryption algorithm from passwords or pass phrases typed by a human will seldom correspond to 128 bits much less 256. To even approach 128--bit equivalence in a pass phrase, at least 10 typical passwords of the kind frequently used in day-to-day work are needed. Weak keys can be somewhat strengthened by special techniques by adding computationally intensive steps which increase the amount of computation necessary to break it. The risks of incorrect usage, implementation and weak keys are in no way unique for AES; these are shared by all encryption algorithms. Provided that the implementation is correct, the security provided reduces to a relatively simple question about how many bits the chosen key, password or pass phrase really corresponds to. Unfortunately this estimate is somewhat difficult to calculate, when the key is not generated by a true random generator.

2.4 Security is relative:

Security is not an absolute; it's a relation between time and cost. Any question about the security of encryption should be posed in terms of how long time, and how high cost will it take an attacker to find a key? Currently, there are speculations that military intelligence services possibly have the technical and economic means to attack keys equivalent to about 90 bits, although no civilian researcher has actually seen or reported of such a capability. Actual and demonstrated systems today, within the bounds of a commercial budget of about 1 million dollars can handle key lengths of about 70 bits. An aggressive estimate on the rate of technological progress is to assume that technology will double the speed of computing devices every year at an unchanged cost.

If correct, 128-bit keys would be in theory be in range of a military budget within 30-40 years. An illustration of the current status for AES is given by the following example, where we assume an attacker with the capability to build or purchase a system that tries keys at the rate of one billion keys per second. This is at least 1 000 times faster than the fasted personal computer in 2004. Under this assumption, the attacker will need about 10 000 000 000 000 000 000 years to try all possible keys for the weakest version, AES-128.

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The key length should thus be chosen after deciding for how long security is required, and what the cost must be to brute force a secret key. In some military circumstances a few hours or days security is sufficient – after that the war or the mission is completed and the information uninteresting and without value. In other cases a lifetime may not be long enough.

3. IMPLEMENTATION:

Advanced Encryption Standard (AES) is a specification for the encryption of electronic data. It has been adopted by the U.S. government and is now used worldwide. The algorithm described by AES is a symmetric-key algorithm, meaning the same key is used for both encrypting and decrypting the data. In the United States, AES was announced by National Institute of Standards and Technology (NIST) as U.S. FIPS PUB 197 (FIPS 197) on November 26, 2001 after a five-year standardization process in which fifteen competing designs were presented and evaluated before it was selected as the most suitable (see Advanced Encryption Standard process for more details). It became effective as a Federal government standard on May 26, 2002 after approval by the Secretary of Commerce.

It is available in many different encryption packages. AES is the first publicly accessible and open cipher approved by the National Security Agency (NSA) for top secret information (see Security of AES, below). Originally called Rijndael, the cipher was developed bytwo Belgian cryptographers, Joan Daemen and Vincent Rijmen, and submitted by them to the AES selection processtwo Belgian cryptographers, Joan Daemen and Vincent Rijmen, and submitted by them to the AES selection process. The name Rijndael (Dutch pronunciation: ['reinda:l][5]) is a play on the names of the two inventors.Strictly speaking, AES is the name of the standard, and the algorithm described is a (restricted) variant of Rijndael. However, in practice the algorithm is also referred to as "AES" (a case of totum pro parte).

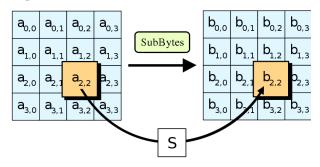
3.2 Description of the cipher:

AES is based on a design principle known as a substitution-permutation network. It is fast in both software and hardware.[6] Unlike its predecessor, DES, AES does not use a Feistel network.

AES has a fixed block size of 128 bits and a key size of 128, 192, or 256 bits, whereas Rijndael can be specified with block and key sizes in any multiple of 32 bits, with a minimum of 128 bits. The blocksize has a maximum of 256 bits, but the keysize has no theoretical maximum. AES operates on a 4×4 column-major order matrix of bytes, termed the state (versions of Rijndael with a larger block size have additional columns in the state). Most AES calculations are done in a special finite field. The AES cipher is specified as a number of repetitions of transformation rounds that convert the input plaintext into the final output of ciphertext. Each round consists of several processing steps, including one that depends on the encryption key. A set of reverse rounds are applied to transform ciphertext back into the original plaintext using the same encryption kev.

4. DISCUSSION: The SubBytes step:

In the SubBytes step, each byte in the matrix is updated using an 8-bit substitution box, the Rijndael S-box. This operation provides the non-linearity in the cipher. The Sbox used is derived from the multiplicative inverse over GF(28), known to have good non-linearity properties. To avoid attacks based on simple algebraic properties, the S-box is constructed by combining the inverse function with an invertible affine transformation. The S-box is also chosen The S-box is also chosen to avoid any fixed points (and so is a derangement), and also any opposite fixed points.



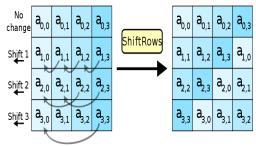
The ShiftRows step:

The ShiftRows step operates on the rows of the state; it cyclically shifts the bytes in each row by a certain offset. For AES, the first ro33w is left unchanged. Each byte of the second row is shifted one to the left. Similarly, the third and fourth rows are shifted by offsets of two and three respectively.



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For the block of size 128 bits and 192 bits the shifting pattern is the same. In this way, each column of the output state of the ShiftRows step is composed of bytes from each column of the input state. (Rijndael variants with a larger block size have slightly different offsets). In the case of the 256-bit block, the first row is unchanged and the shifting for second, third and fourth row is 1 byte, 3 bytes and 4 bytes respectively—this change only applies for the Rijndael cipher when used with a 256-bit block, as AES does not use 256-bit blocks.



Brief Description of Rijndael Algorithm:

Rijndael algorithm consists of encryption, decryption and key schedule algorithm. The main operations of the encryption algorithm among the three parts of Rijndael algorithm include: bytes substitution (SubBytes), the row shift (ShiftRows), column mixing (MixColumns), and the round key adding (AddRoundKey). It is shown as Fig. 1.

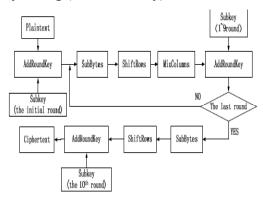


Figure. The structure of Rijndael encryption algorithm

5. Concluding Remarks :

AES-128 algorithm for encryption and decryption is implemented in Virtex-5 FPGA. With the designing of all the operations as LUTs and ROMs, the proposed architectureachieves a throughput of 3.74 Gbps and thereby utilizing only 1% of slices in the targeted FPGA. Since the speed is higher

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than the already reported systems, hence the proposed design serves as the best high speed encryption algorithm and is thus suitable for various applications. Moreover with less area utilization, the proposed design can be embedded with other larger designs as well.

6. Experimental Results:

Simulation For Aes

			2,275,014 ps						
Name	Yalue	2,274,900 ps 2,	75,000 ps	2,275,100 ps	2,275,200 ps	2,275,300 ps	2,275,400 ps	2,275,500 ps	2,275
l <mark>la</mark> dk	1								
🕨 🕌 dətain(127:0)	10101011111	101010111111111111	101010101010	1.1111111111111111	01010101010101	1111111111110000		01010101010100	000
▶ <table-of-contents> key(127:0]</table-of-contents>	000000000000		01001001001001		000000111111111	11111111111111111		1010101010101010	1111
▶ <table-of-contents> dataout[127:0]</table-of-contents>	10010010000	1001001000011101010	0011001110010	00010001110011110	00001101101111001	0000110001110001	1101110000111000	0100110000001110	00101:
🕨 💐 r1_cut(127:0)	11101001011	110100101110100011	0100100111001	01111111101010001	0101000010101110	0111101010001111	1100101100100000	0010000000000000	00001:
▶ 👹 r2_cut[127:0]	10010010100	1001001010000001010	0110011100111	10100011101110001	1000001100100100	0101001010001110	0100010100000111	1000110010010100	10100
▶ 👹 r3_cut[127:0]	00101000111	0010100011101011110	1000111100111	0011110001010001	1111111111001001	0011010011010011	100001100000011	0110011110011011	11100
▶ 👹 r4_cut[127:0]	01010011011	0101001101110111110	1101000111111	01110001011101011	011011100000010	0001101101001100	1110101111101110	0111000000011001	00111
▶ 👹 r5_cut[127:0]	10010001011	1001000101101011111	0101010001111	10100110010111101	1100101100010100	0011000010100100	1010101001111101	0111010011110110	00110:
▶ 👹 r6_cut[127:0]	00010100010	0001010001000101111	0000000100011	0001110111110011	110101010101010110	011011001111101	110110010000011	01001100110001000	01111
▶ 👹 r7_cut[127:0]	00011011101	0001101110101010001	1011101000100:	1001011010001000	000101101111101	000111010101010101	1000000100111001	0011100010110011	01011
▶ 👹 r8_out[127:0]	11010100010	101010001011100001	0100111000000	1000011111000111	01001011011110111	0001100100010000	0101111110000101	0000110100011101	1010
▶ 👹 r9_out[127:0]	10011001001	1001100100101101100	1011001100011	01011110101110100	00001000100001000	101001000111101	1010110000101000	0010101010010111	10111
🕨 👹 keyauti (127:0)	00010111000	0001011100010110000	0110000101100	0101110001011000	01011001101001111	0100011101001111	1001100101100001	01110001011000010	11001:
▶ 👹 keyaut2[127:0]	00010100000	0001010000010110000	0110000101100	0101000001011000	01011001101001111	0101111101001111	1001100101100001	01000001011000010	11001:
🕨 🙀 keyaut3[127:0]	00010010000	0001001000010110000	0110000101100	0100100001011000	01011001101001111	0110111101001111	1001100101100001	0100001011000010	11001:
🕨 🙀 keyaut4(127:0)	00011110000	0001111000010110000	0110000101100	0111100001011000	01011001101001111	0000111101001111	1001100101100001	1100001011000010	11001:
🕨 🙀 keyauti(127:0)	00000110000	0000011000010110000	0110000101100	00001100001011000	01011001101001111	100111101001111	1001100101100000	01100001011000010	11001:

Simulation For Rounds

				1,0	oo'tos hel					
Name	¥alue		1,030,000 ps	Ţ	1,030,200 ps	1,030,400 ps	1,030,600 ps	1,030,800 ps	1,031,000 ps	1,031,200 ps
🕨 <table-of-contents> rd(3:0)</table-of-contents>	1001						1001			
埍 ck	0									
🕨 <table-of-contents> datain(127:0]</table-of-contents>	1111111111	1111	11111101010101	010	110010110000000	0000011010101010101	0000000000000000	000000111111111	10000000000000	
🕨 <table-of-contents> keyin(127:0)</table-of-contents>	01010101010	01010	010101010101010	01	010101010101010101	011111111111111111	нынынын		00000000000000000	0101010101010101010
▶ <table-of-contents> keyout[127:0]</table-of-contents>	00000000000	0000	00000110110001	01	1000110110010101	01110010011100100	110010011010101	0011011000110110	100100110101010	01101100011011011
▶ 👹 b00_sb(7:0)	00010110						00010110			
▶ 👹 b01_sb(7:0)	01100011						01100011			
▶ 👹 b02_sb(7:0]	01100011						01100011			
▶ 👹 b03_sb(7:0]	00010110						00010110			
▶ 👹 b10_sb(7:0]	10000111						10000111			
▶ 👹 b11_sb(7:0)	10100010						10100010			
▶ 👹 b12_st(7:0]	01111011						01111011			
▶ 👹 b13_sb(7:0]	00010110						00010110			
▶ 👹 b20_sb(7:0]	01100010						01100010			
▶ 👹 b21_sb(7:0]	10101100						10101100			
▶ 👹 b22_sb(7:0)	00010110						00010110			
▶ 💐 b23_sb(7:0]	00010110						00010110			
▶ 👹 b30_sb(7:0)	01110001						01110001			



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Simulation waveform for encryption

Value	0 ns 500 ns
01234567	XXXX
1	
76bde35a	(0000) 76bde35a4772657f2bc56e1afd456767
00000000	(0000X 000000000000000000000000000000
	01234567 1 76bde35a

Simulation waveform for Decryption

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